# Non-Hamiltonian and Non-Traceable Regular 3-Connected Planar Graphs

Nico Van Cleemput Carol T. Zamfirescu

Combinatorial Algorithms and Algorithmic Graph Theory
Department of Applied Mathematics, Computer Science and Statistics
Ghent University





Introduction

Definitions

Cubic

Quartic

Quintic

Summary

2 Cubic

Essentially 4-connected

3 Quartic

Upper bound c<sub>4</sub>

Lower bound c₄

Upper bound p4

Lower bound p<sub>4</sub>

4 Quintic Upper bound  $p_5$ 

5 Conclusion Summary





- Here, a polyhedron is a planar 3-connected graph.
- The word "regular" is used exclusively in the graph-theoretical sense of having all vertices of the same degree.
- By Euler's formula, there are *k*-regular polyhedra for exactly three values of *k*: 3, 4, or 5.



- Let c<sub>k</sub> be the order of the smallest non-hamiltonian k-regular polyhedron.
- Let  $p_k$  be the order of the smallest **non-traceable** k-regular polyhedron.



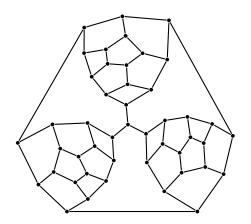
#### Cubic polyhedra – hamiltonicity

- Tait conjectured in 1884 that every cubic polyhedron is hamiltonian.
- The conjecture became famous because it implied the Four Colour Theorem (at that time still the Four Colour Problem)



Introduction Cubic Quartic Quintic Conclusion Definitions Cubic Quartic Quintic Summary

#### Cubic polyhedra – hamiltonicity

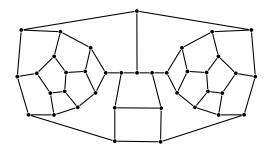


The first to construct a counterexample (of order 46) was Tutte in 1946



Introduction Cubic Quartic Quintic Conclusion Definitions Cubic Quartic Quintic Summary

#### Cubic polyhedra – hamiltonicity



Lederberg, Bosák, and Barnette (pairwise independently) described a smaller counterexample having 38 vertices.



## Cubic polyhedra – hamiltonicity

After a long series of papers by various authors (e.g., Butler, Barnette, Wegner, Okamura), Holton and McKay showed that all cubic polyhedra on up to 36 vertices are hamiltonian.

Theorem (Holton and McKay, 1988)

 $c_3 = 38$ 



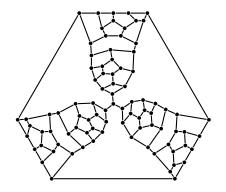
#### Cubic polyhedra – traceability

- Balinski asked whether cubic non-traceable polyhedra exist
- Brown and independently Grünbaum and Motzkin proved the existence of such graphs
- Klee asked for determining p<sub>3</sub>



Introduction Cubic Quartic Quintic Conclusion Definitions Cubic Quartic Quintic Summary

#### Cubic polyhedra – traceability



In 1970 T. Zamfirescu constructed this cubic non-traceable planar graph on 88 vertices



#### Based on work of Okamura, Knorr improved a result of Hoffmann by showing that all cubic planar graphs on up to 52 vertices are traceable.

Theorem (Knorr, 2010 and Zamfirescu, 1970)

$$54 \le p_3 \le 88$$



## Quartic polyhedra – hamiltonicity

- Following work of Sachs from 1967 and Walther from 1969, Zaks proved in 1976 that there exists a quartic non-hamiltonian polyhedron of order 209.
- The actual number given in Zaks' paper is false, as pointed out in work of Owens — therein the correct number can be found.



#### Quartic polyhedra – hamiltonicity

#### Theorem (Sachs, 1967)

If there exists a non-hamiltonian (non-traceable) cubic polyhedron of order n, then there exists a non-traceable (non-hamiltonian) quartic polyhedron on  $\frac{9n}{2}$  vertices.

On page 132 of Bosák's book it is claimed that converting the Lederberg-Bosák-Barnette graph with this method gives a quartic non-hamiltonian polyhedron of order 161. However, the correct number should be  $38 \times \frac{9}{2} = 171$ .

Theorem (Sachs, 1967 combined with Bosák, 1990)

 $c_4 < 171$ 



## Quartic polyhedra – traceability

- **Zaks showed that**  $p_4 \le 484$
- Using Sachs' theorem on Zamfirescu's 88-vertex graph gives a non-traceable quartic polyhedron on 396 vertices.

Theorem (Sachs, 1967 combined with Zamfirescu, 1970)

 $p_4 \le 396$ 



## Quintic polyhedra

- Previous work includes papers by Walther, as well as Harant, Owens, Tkáč, and Walther.
- **Zaks** showed that  $c_5 \le 532$  and  $p_5 \le 1232$ .
- Owens proved that  $c_5 \le 76$  and  $p_5 \le 128$ .

#### Theorem (Owens, 1980)

$$c_5 \leq 76$$

$$p_5 \le 128$$



Introduction Cubic Quartic Quintic Conclusion Definitions Cubic Quartic Quintic Summary

## Summary

|         | Hamiltonicity               | Traceability                    |
|---------|-----------------------------|---------------------------------|
| Cubic   | <i>c</i> <sub>3</sub> = 38  | 54 ≤ <i>p</i> <sub>3</sub> ≤ 88 |
| Quartic | <i>c</i> <sub>4</sub> ≤ 171 | $p_4 \leq 396$                  |
| Quintic | <i>c</i> <sub>5</sub> ≤ 76  | <i>p</i> <sub>5</sub> ≤ 128     |



# Cubic Polyhedra



## Essentially 4-connected cubic polyhedra

Theorem (Aldred, Bau, Holton, and McKay, 2000)

Every essentially 4-connected cubic planar graph of order at most 40 is hamiltonian. Furthermore, there exist non-hamiltonian examples of order 42.

Theorem (Van Cleemput and Zamfirescu, 2018)

There exists a non-hamiltonian essentially 4-connected cubic polyhedron of order n if and only if n is even and n > 42.



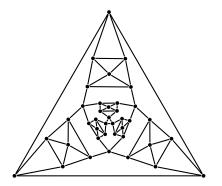
#### Quartic Polyhedra



#### Upper bound hamiltonicity

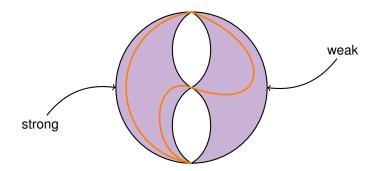
#### Theorem (Van Cleemput and Zamfirescu, 2018)

 $c_4 \le 39$ 

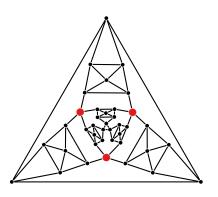


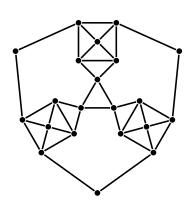


#### Upper bound hamiltonicity



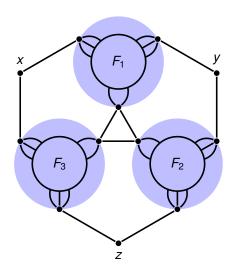






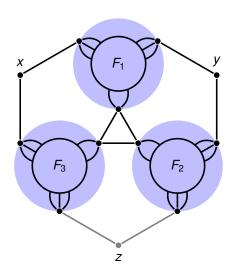


#### Upper bound hamiltonicity





#### Upper bound hamiltonicity



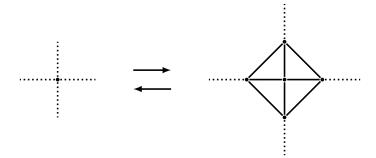


#### Lower bound hamiltonicity

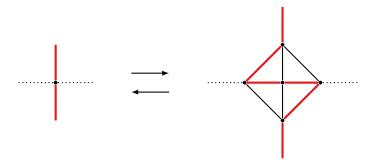
Check all quartic polyhedra for being hamiltonian.

Simple backtracking algorithm that tries to construct a cycle from some vertex.

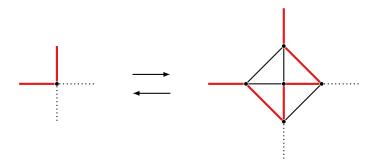




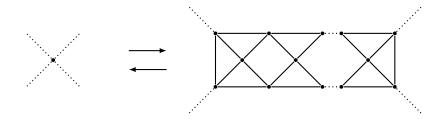






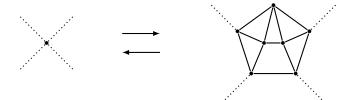








Introduction Cubic Quartic Quintic Conclusion Upper bound  $c_4$  Lower bound  $c_4$  Upper bound  $c_5$  Upper bound  $c_6$  Upper bound  $c_8$  Upper bound  $c_8$  Upper bound  $c_9$  U





#### Lower bound hamiltonicity

#### Theorem (Van Cleemput and Zamfirescu, 2018)

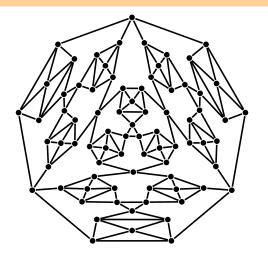
 $c_4 \ge 35$ 

| Vertices | Time         |
|----------|--------------|
| 25       | 9.6 minutes  |
| 26       | 42.1 minutes |
| 27       | 3.2 hours    |
| 28       | 15.1 hours   |
| 29       | 3.1 days     |
| 30       | 15.3 days    |
| 31       | 78.2 days    |
| 32       | 1.1 years    |
| 33       | 5.9 years    |
| 34       | 37.9 years   |
|          |              |

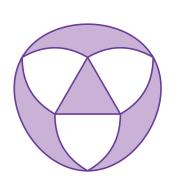


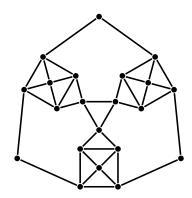
#### Theorem (Van Cleemput and Zamfirescu, 2018)

 $p_4 \le 78$ 









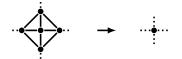
$$21 \times 4 - 6 = 78$$
 vertices

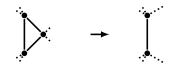


# Lower bound traceability

#### Lemma (Van Cleemput and Zamfirescu, 2018)

$$p_4 \ge c_4 + 1$$



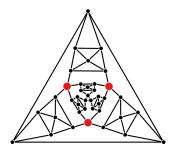


#### Theorem (Van Cleemput and Zamfirescu, 2018)

 $p_4 \ge 36$ 

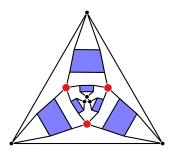


## Further properties



- Not homogeneously traceable
- Circumference is  $5 \times 6 + 4 = 34$ .

## Further properties



- For each  $n \ge 39$  there is a quartic polyhedron on n vertices that is not homogeneously traceable.
- For the family  $\mathcal{G}$  of quartic polyhedra, the shortness coefficient  $\rho(\mathcal{G})$  is at most  $\frac{5}{6}$ :

$$\rho\left(\mathcal{G}\right) = \liminf_{G \in \mathcal{G}} \frac{\mathsf{circ}(G)}{|V(G)|} \leq \liminf_{k \to \infty} \frac{5k+4}{6k+3} = \frac{5}{6}$$

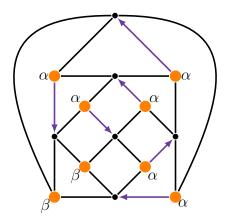
# Quintic Polyhedra

Upper bound p5

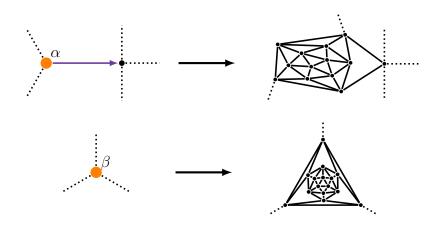


#### Theorem (Van Cleemput and Zamfirescu, 2018)

 $p_5 \le 108$ 







$$6\times13+2\times15=108$$
 vertices



Theorem (Van Cleemput and Zamfirescu, 2018)

There exists a quintic non-traceable polyhedron of order n for every even  $n \ge 108$ .



## Summary

|         | Hamiltonicity  | Traceability   |
|---------|--|--|
| Cubic   | $c_3 = 38$   | $54 \le p_3 \le 88$  |
| Quartic | $\begin{array}{c} c_4 \leq 171 \\ 35 \leq c_4 \leq 39 \end{array}$ | _ <del>p<sub>4</sub>≤ 39</del> 6<br>36 ≤ p <sub>4</sub> ≤ 78 |
| Quintic | $c_5 \leq 76$ $38 \leq c_5 \leq 76$                                | $p_5 \le 128$ $38 \le p_5 \le 108$                           |



- Increase lower bound for c<sub>4</sub>
  - number of 3-cuts
  - required subgraphs
- Lower bounds for quintic case
- Lower bounds for traceability
- Upper bound for hamiltonicity of quintic case

